

# Intro to Graphs

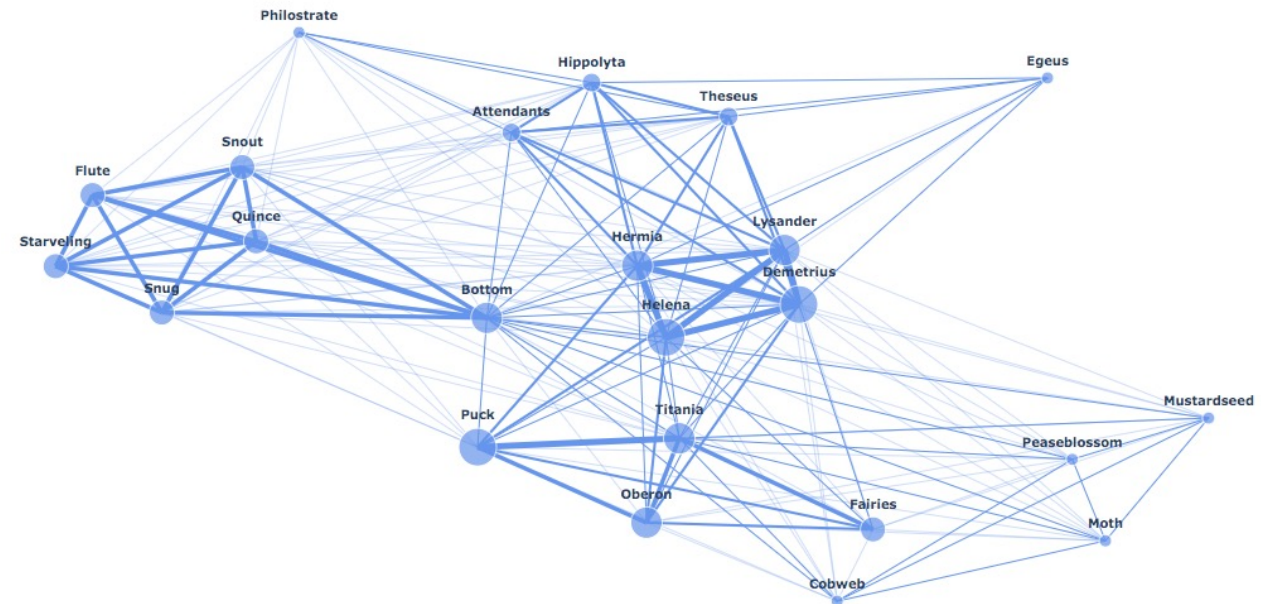
## Theoretical foundations

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Giuseppe Averta

Carlo Masone

Francesca Pistilli





Introduction to Graphs

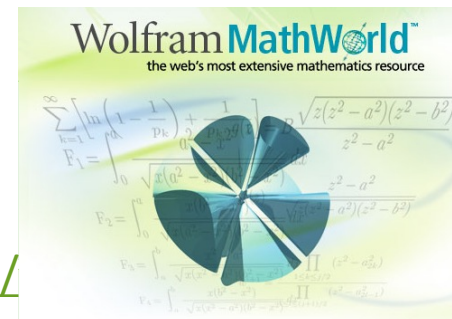


# DEFINITION: GRAPH

# Definition: Graph

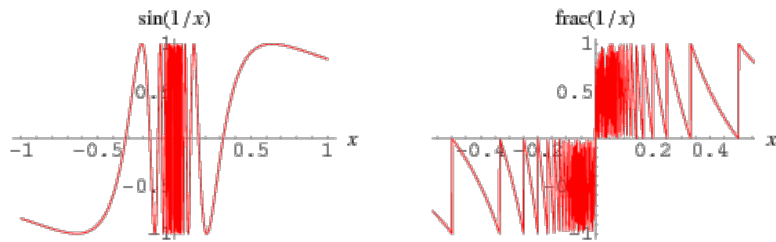
- A **graph** is a collection of **points** and **lines** connecting some (eventually empty) subset of them.
- The points of a graph are typically known as **graph vertices**, but may also be called “nodes” or simply “points.”
- The lines connecting the vertices of a graph are called **graph edges**, but may also be called “arcs” or “lines.”

<http://mathworld.wolfram.com/>

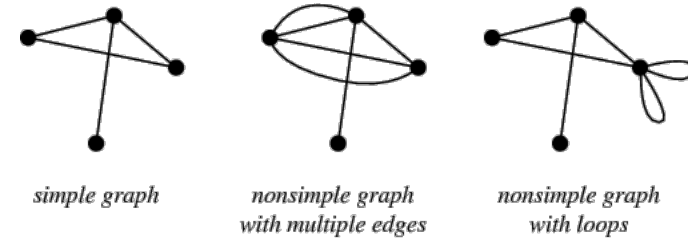


# Big warning: Graph $\neq$ Graph $\neq$ Graph

- Graph (plot)
- (italiano: grafico)

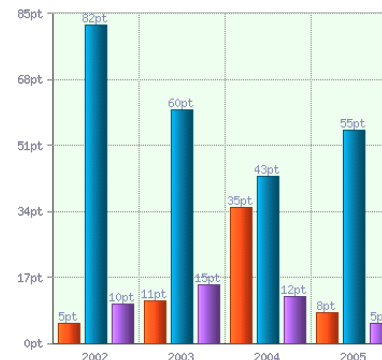


- Graph (maths)
- (*italiano: grafo*)



$\neq$

**Graph (chart)**  
(italiano: grafico)



# History

- The study of graphs is known as **graph theory**, and was first systematically investigated by D. König in the 1930s
- Euler's proof about the *walk across all seven bridges of Königsberg* (1736), now known as the *Königsberg bridge* problem, is a famous precursor to graph theory.
- Indeed, the study of various sorts of paths in graphs has many applications in real-world problems.

# Königsberg Bridge Problem

- Can the 7 bridges of the city of Königsberg over the river Pregel all be traversed in a single trip without doubling back, with the additional requirement that the trip ends in the same place it began?

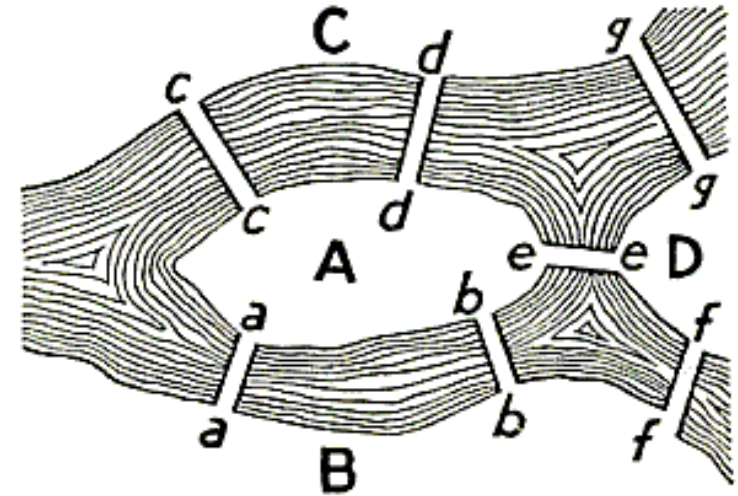
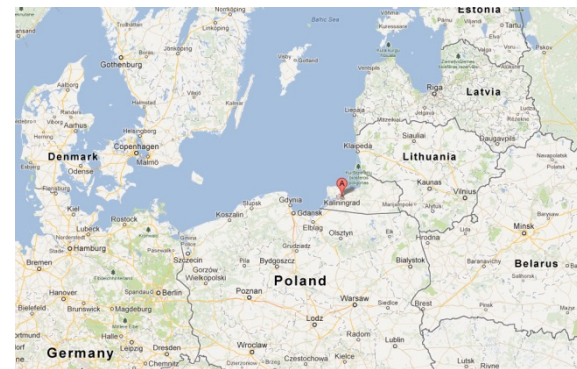


FIGURE 98. *Geographic Map:  
The Königsberg Bridges.*



Today: Kaliningrad, Russia

# Königsberg Bridge Problem

- Can the 7 bridges of the city of Königsberg over the river Pregel all be traversed in a single trip without doubling back, with the requirement that the trip ends at the same place it began?

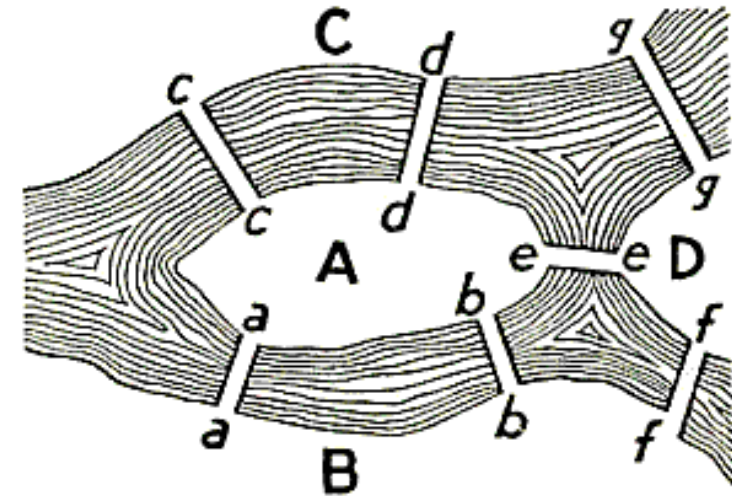
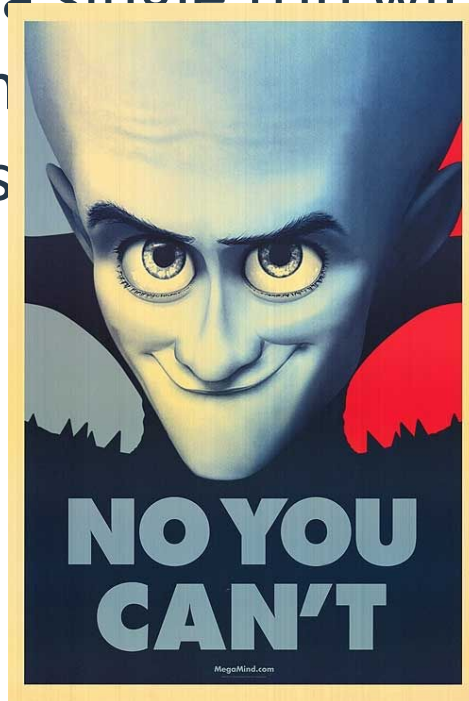
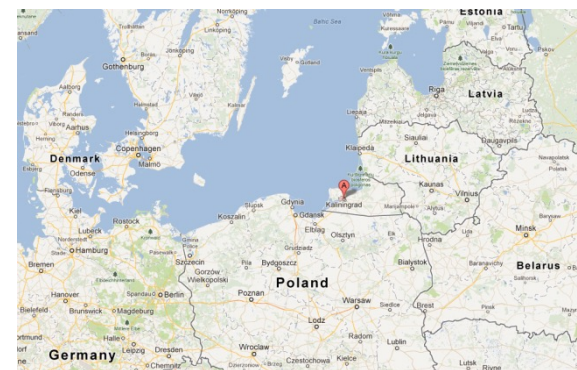


FIGURE 98. *Geographic Map:  
The Königsberg Bridges.*

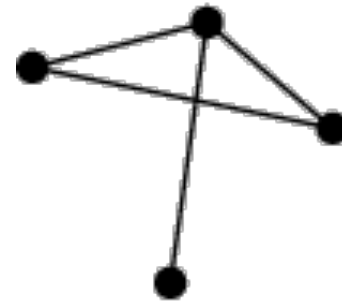


Today: Kaliningrad, Russia

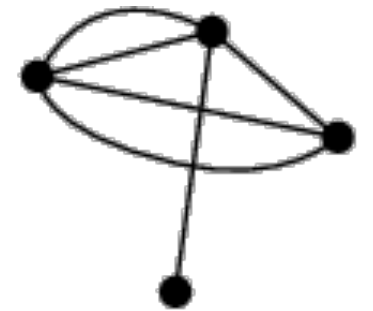


# Types of graphs: edge cardinality

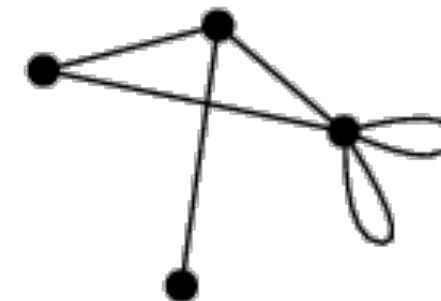
- Simple graph:
  - At most one edge (i.e., either one edge or no edges) may connect any two vertices
- Multigraph:
  - Multiple edges are allowed between vertices
- Loops:
  - Edge between a vertex and itself
- Pseudograph:
  - Multigraph with loops



*simple graph*



*multigraph*

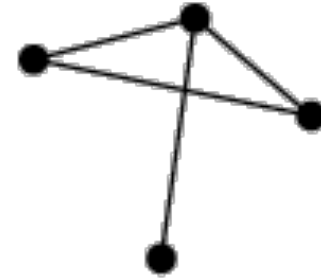


*pseudograph*

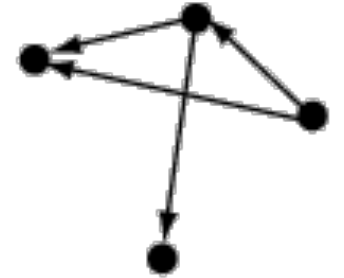


# Types of graphs: edge direction

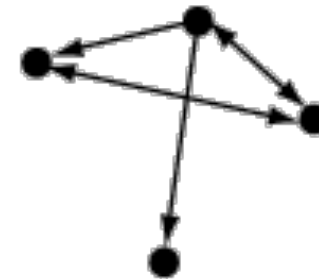
- Undirected
- Oriented
  - Edges have **one** direction (indicated by arrow)
- Directed
  - Edges may have **one or two** directions
- Network
  - Oriented graph with weighted edges



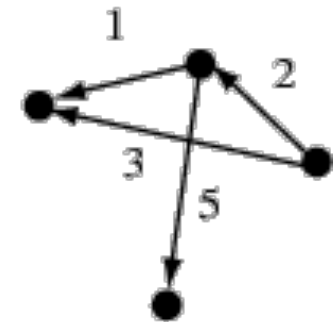
*undirected graph*



*oriented graph*



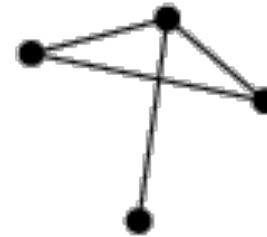
*directed graph*



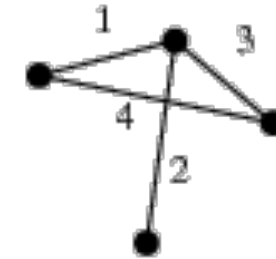
*network*

# Types of graphs: labeling

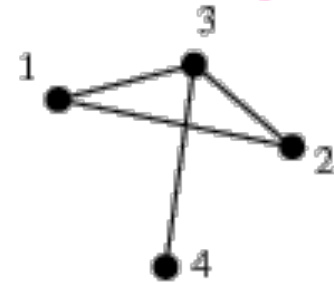
- Labels
  - None
  - On Vertices
  - On Edges
- Groups (=colors)
  - Of Vertices
    - no edge connects two identically colored vertices
  - Of Edges
    - adjacent edges must receive different colors
  - Of both



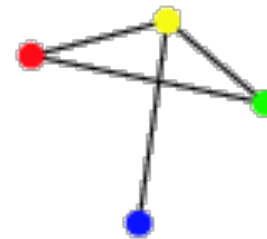
*unlabeled graph*



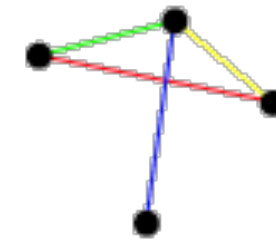
*edge-labeled graph*



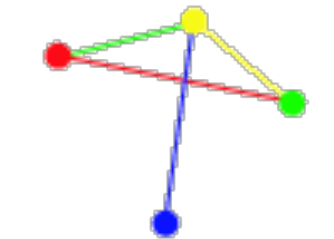
*vertex-labeled graph*



*vertex-colored graph*



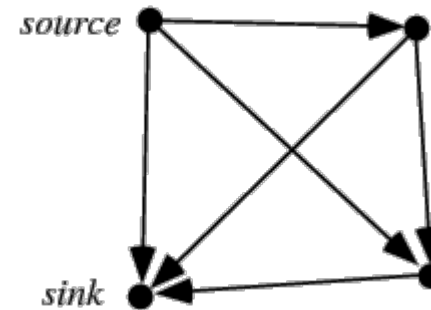
*edge-colored graph*



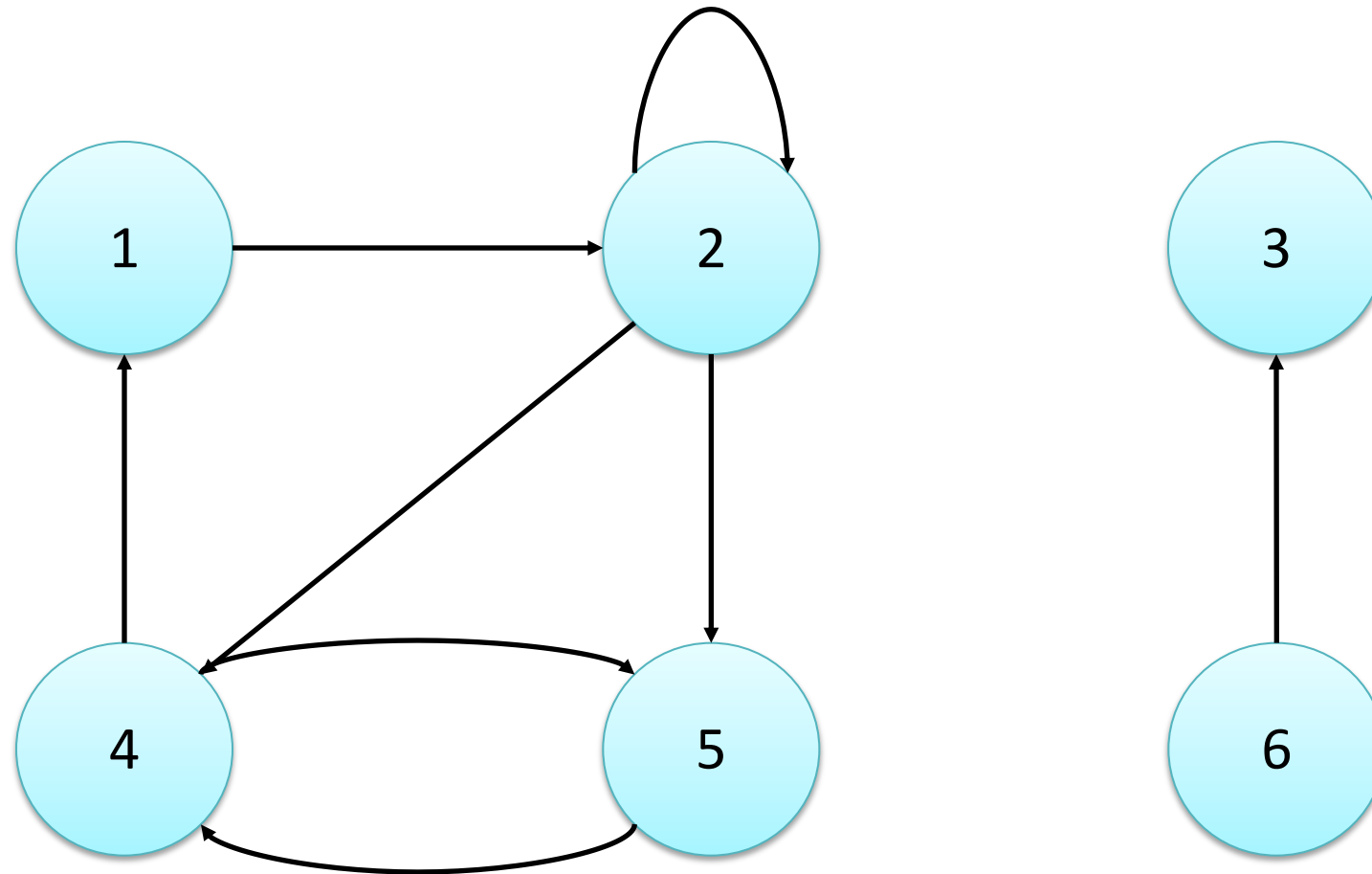
*vertex- and edge-colored graph*

# Directed and Oriented graphs

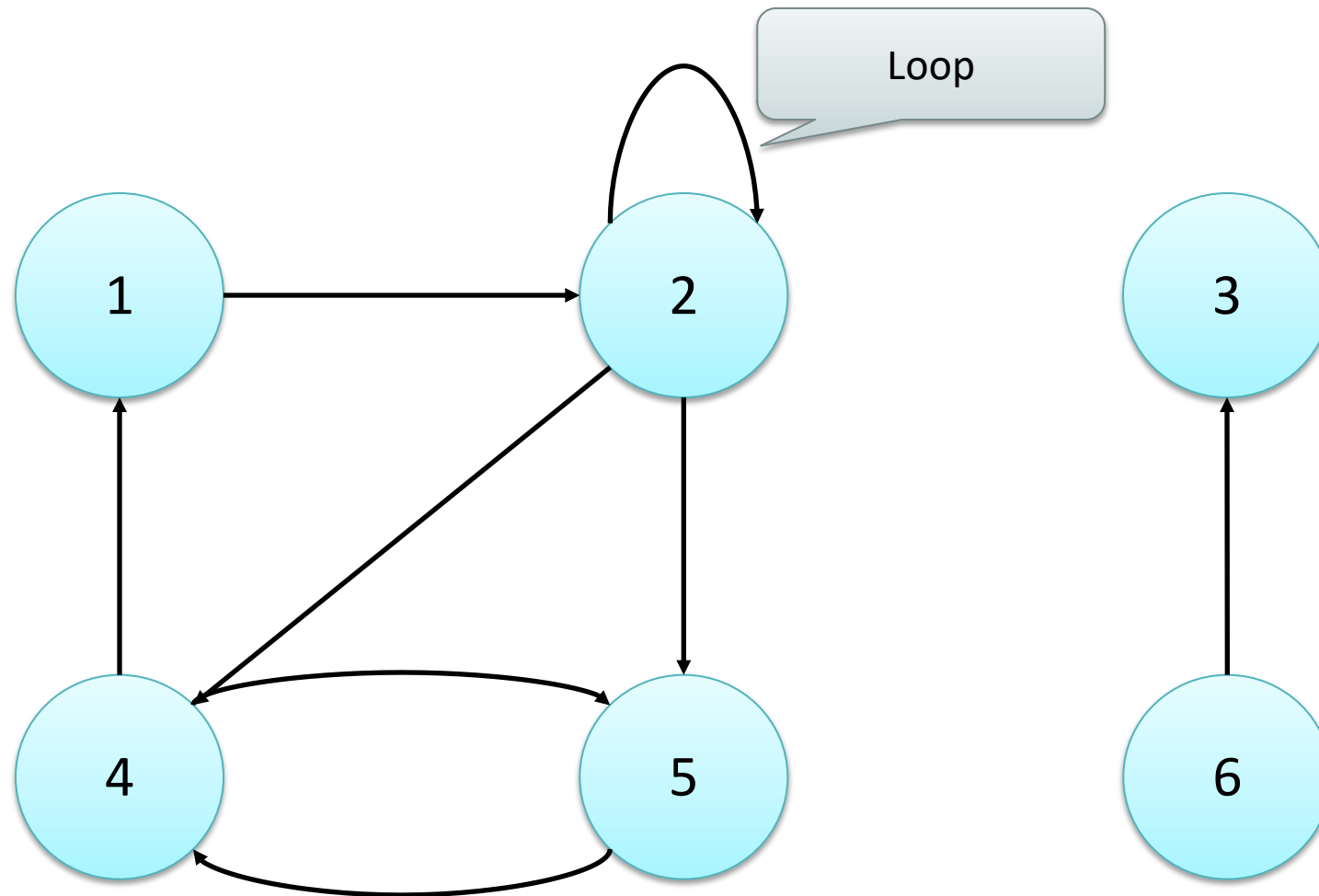
- A Directed Graph (*di-graph*)  $G$  is a pair  $(V,E)$ , where
  - $V$  is a (finite) set of *vertices*
  - $E$  is a (finite) set of *edges*, that identify a binary relationship over  $V$ 
    - $E \subseteq V \times V$



# Example



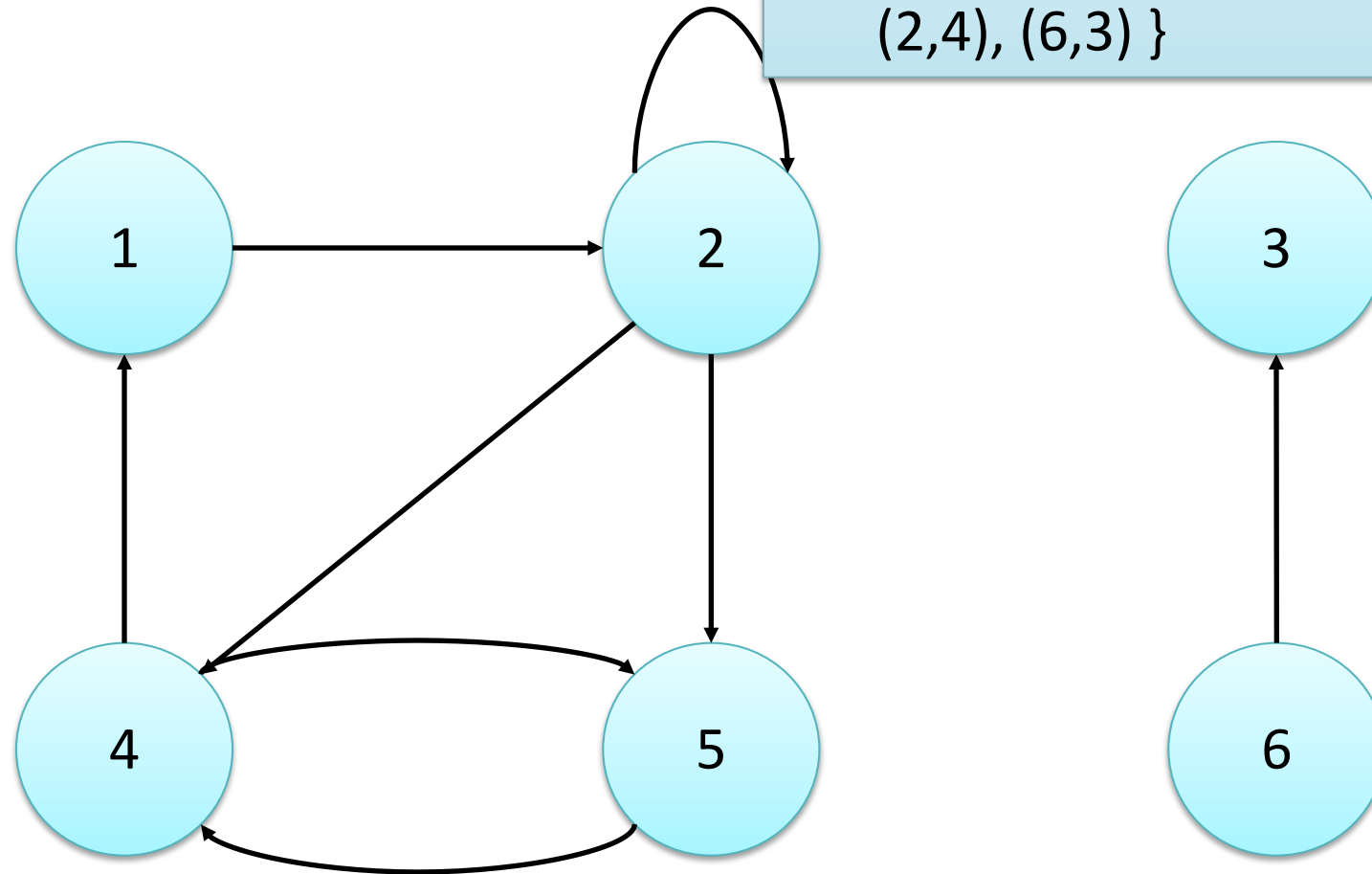
# Example



# Example

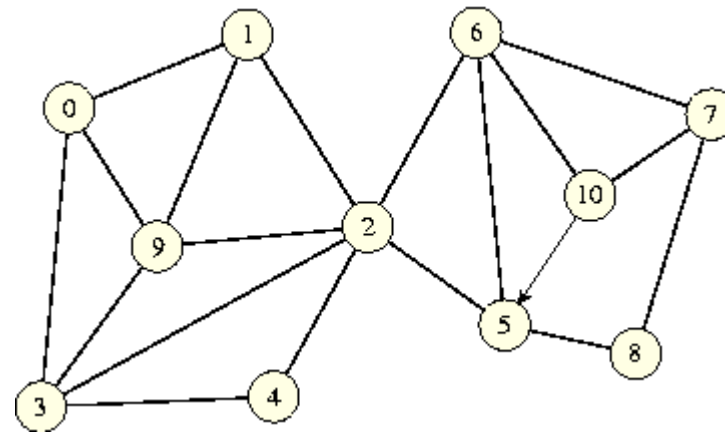
$V = \{1,2,3,4,5,6\}$

$E = \{ (1,2), (2,2), (2,5), (5,4), (4,5), (4,1), (2,4), (6,3) \}$



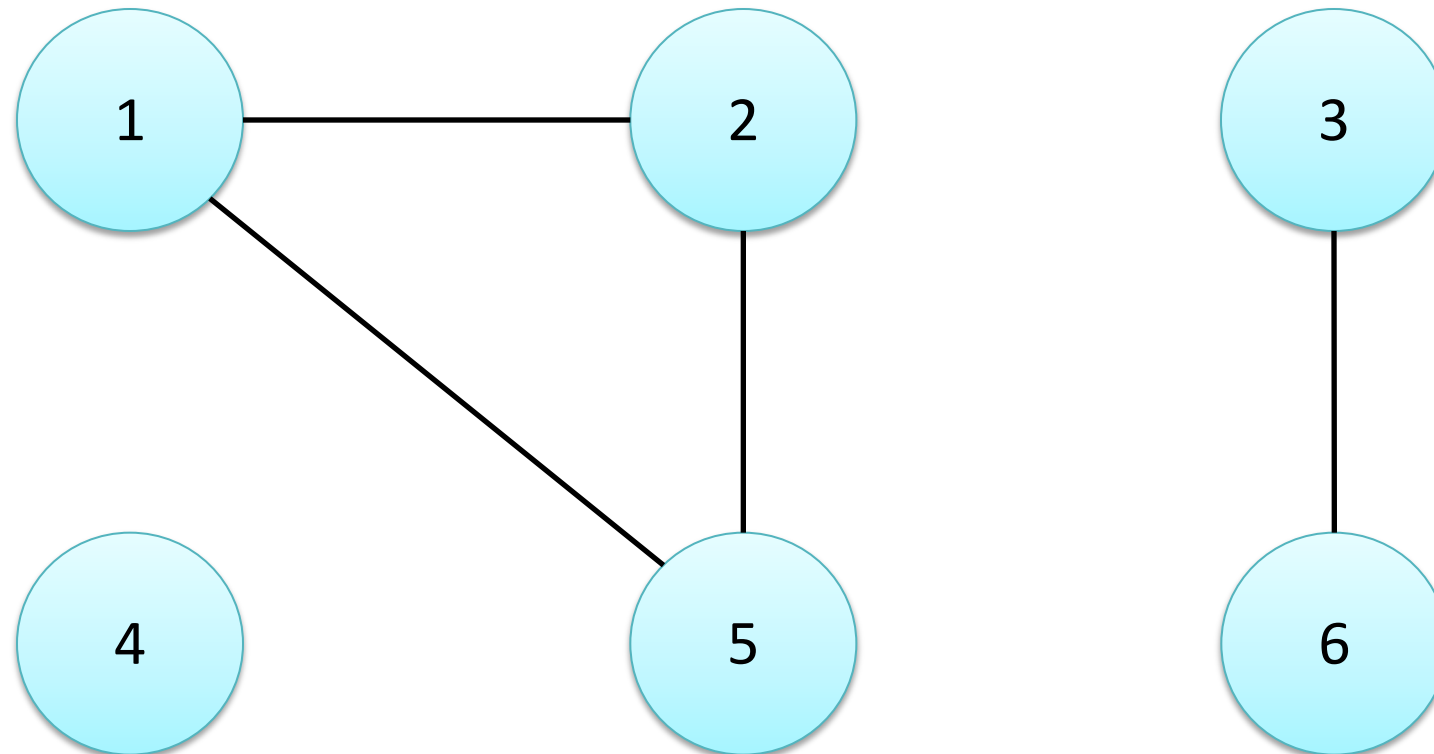
# Undirected graph

- An **Undirected** Graph is still represented as a tuple  $G=(V,E)$ , but the set  $E$  is made of **non-ordered pairs** of vertices

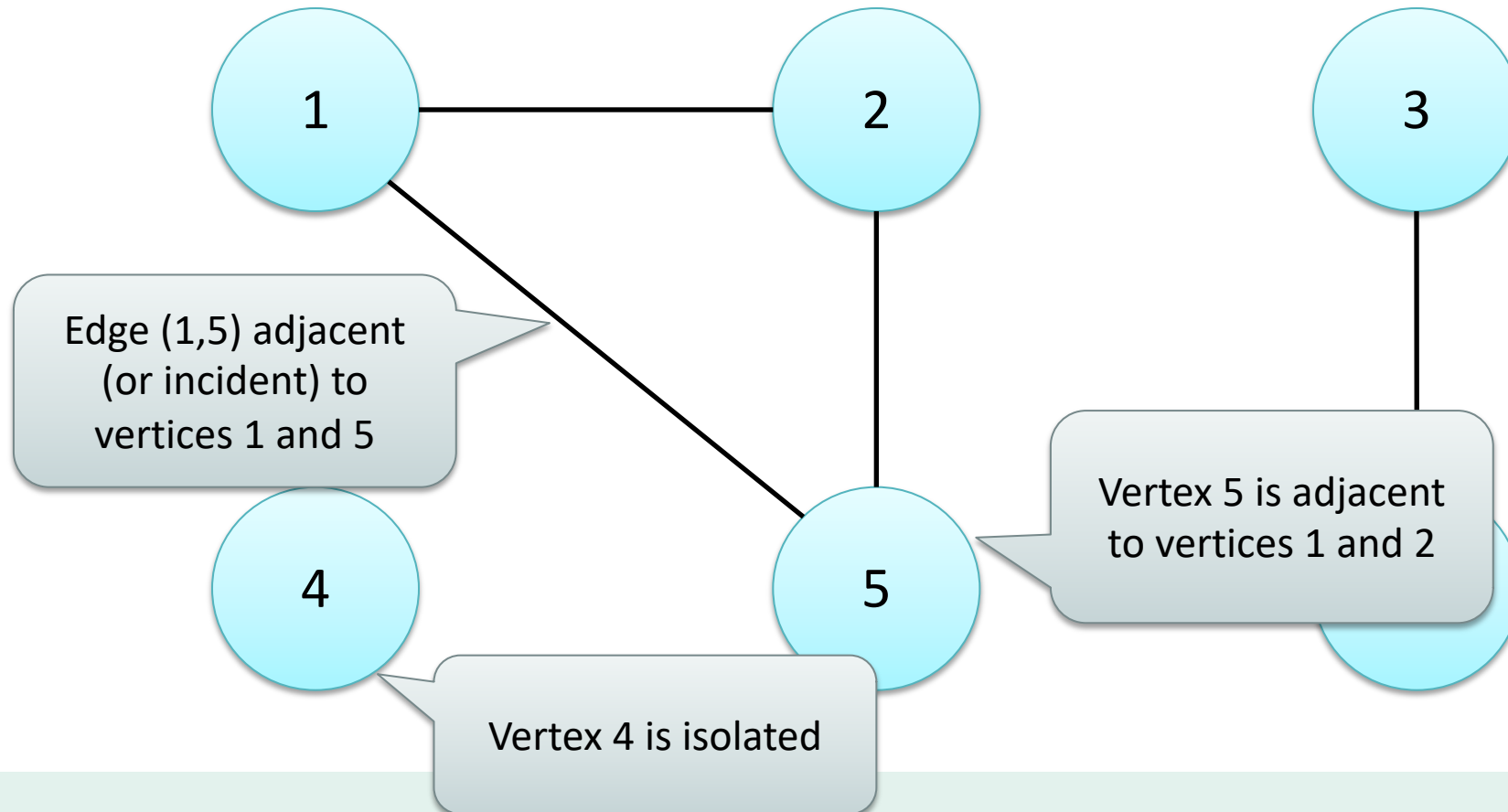




# Example

$$V = \{ 1, 2, 3, 4, 5, 6 \}$$
$$E = \{ \{1,2\}, \{2,5\}, \{5,1\}, \{6,3\} \}$$


# Example

$$V = \{ 1, 2, 3, 4, 5, 6 \}$$
$$E = \{ \{1,2\}, \{2,5\}, \{5,1\}, \{6,3\} \}$$




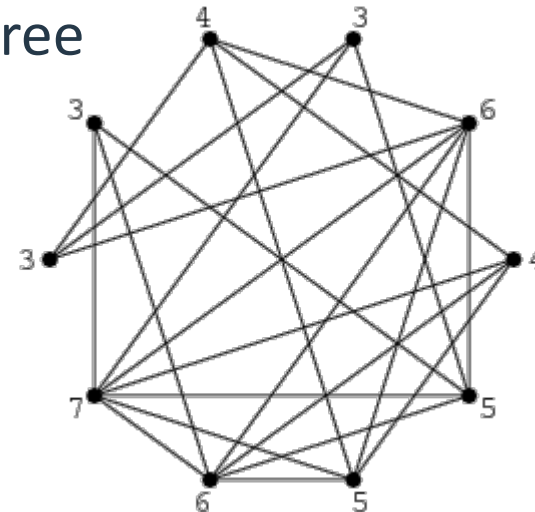
Introduction to Graphs



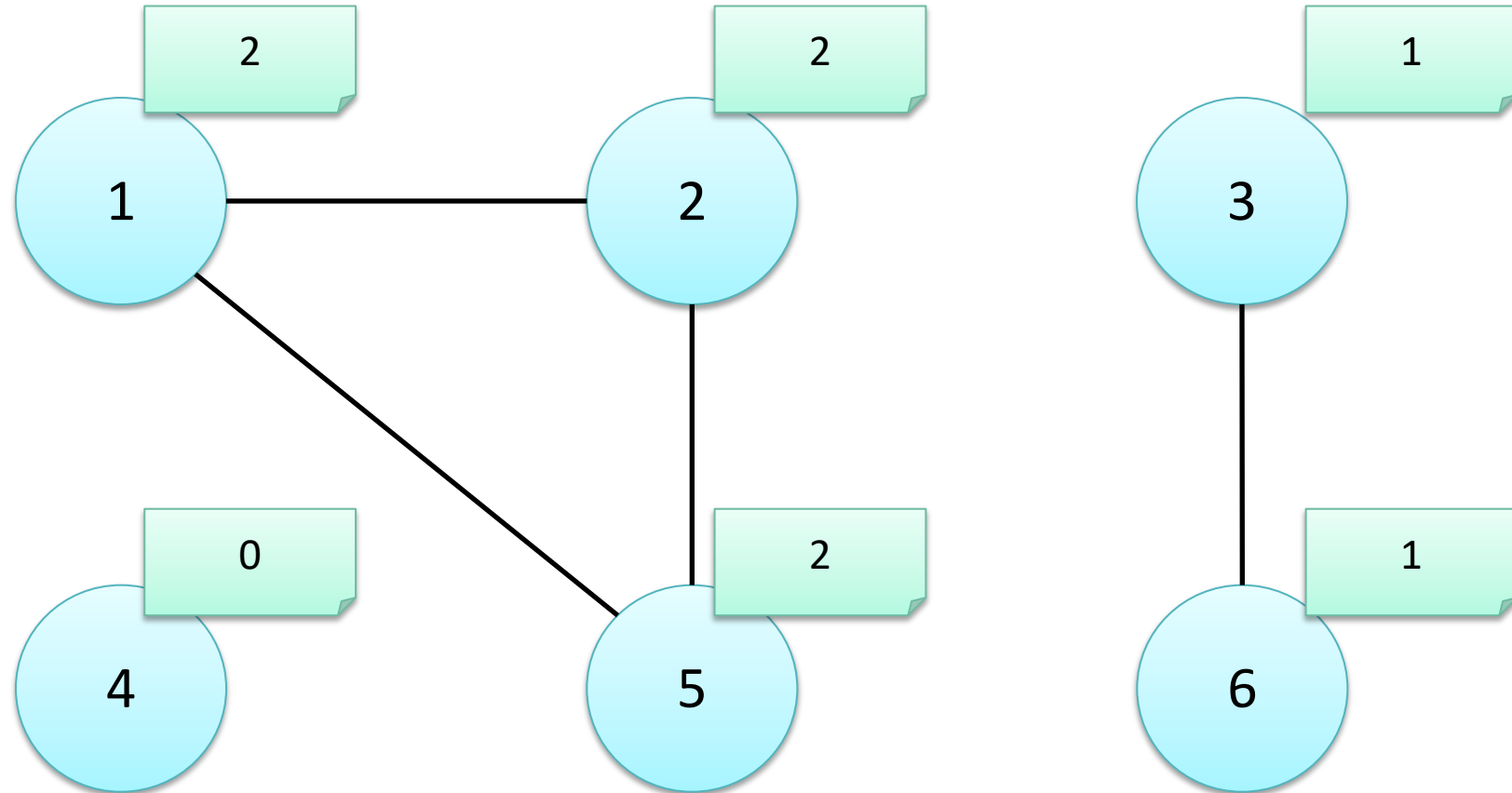
# RELATED DEFINITIONS

# Degree

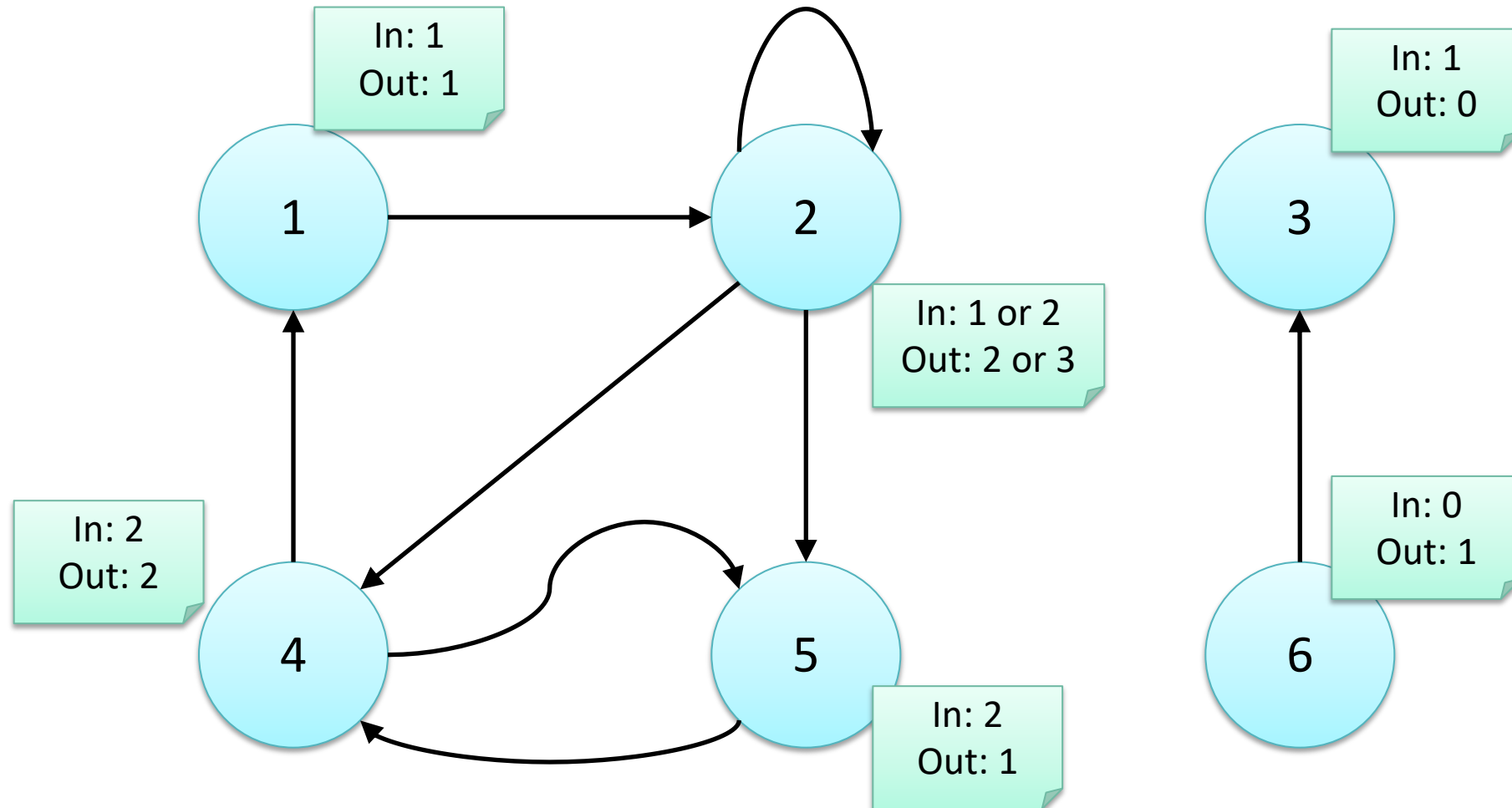
- In an *undirected* graph,
  - the **degree** of a vertex is the number of incident edges
- In a *directed* graph
  - The **in-degree** is the number of incoming edges
  - The **out-degree** is the number of departing edges
  - The **degree** is the sum of in-degree and out-degree
- A vertex with degree 0 is **isolated**



# Degree



# Degree



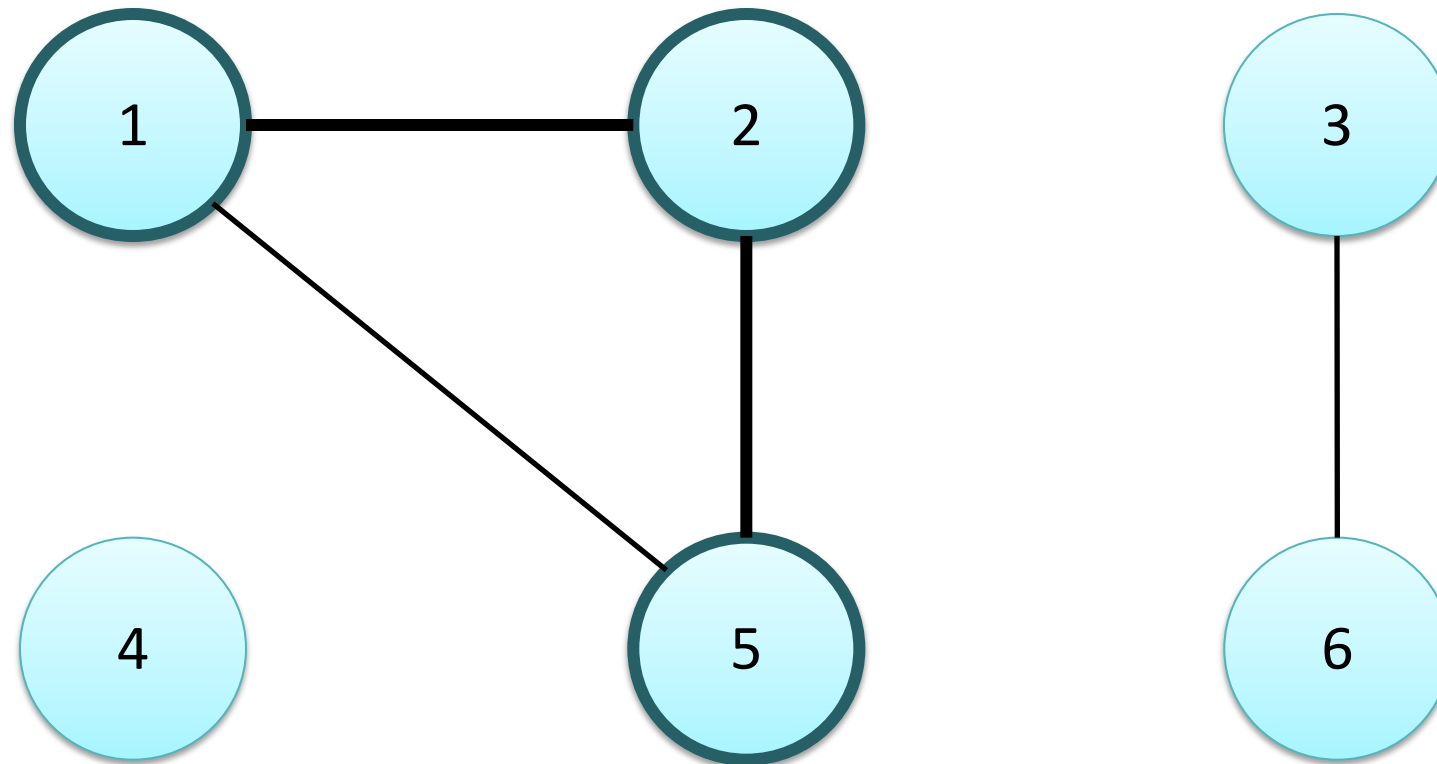
# Paths

- A **path** on a graph  $G=(V,E)$ , also called a trail, is a sequence  $\{v_1, v_2, \dots, v_n\}$  such that:
  - $v_1, \dots, v_n$  are vertices:  $v_i \in V$
  - $(v_1, v_2), (v_2, v_3), \dots, (v_{n-1}, v_n)$  are graph edges:  $(v_{i-1}, v_i) \in E$
  - $v_i$  are distinct (for “simple” paths).
- The **length** of a path is the number of edges  $(n-1)$
- If there exist a path between  $v_A$  and  $v_B$  we say that  $v_B$  is **reachable** from  $v_A$



# Example

Path = ( 1, 2, 5 )  
Length = 2



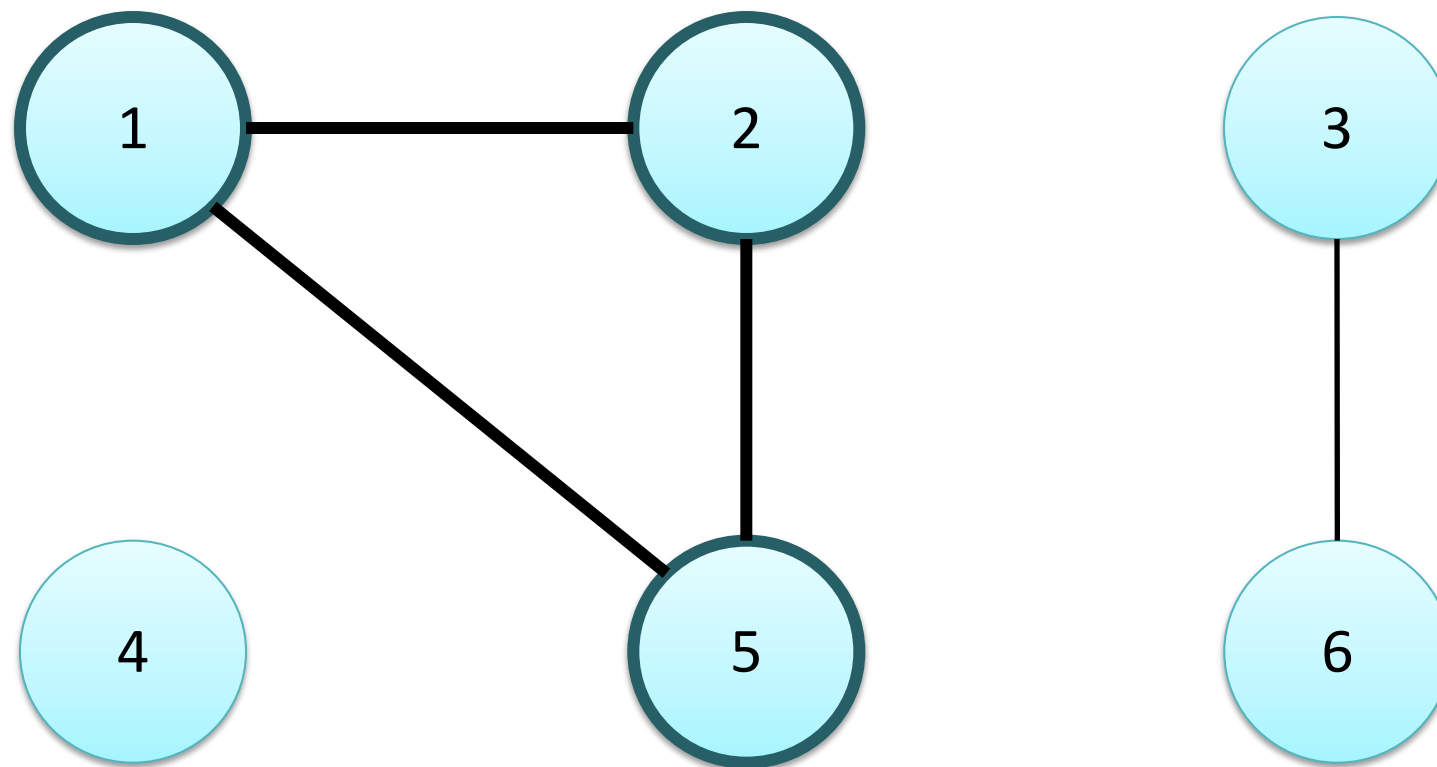


# Cycles

- A cycle is a path where  $v_1 = v_n$
- A graph with no cycles is said acyclic

# Example

Path = ( 1, 2, 5, 1 )  
Length = 3

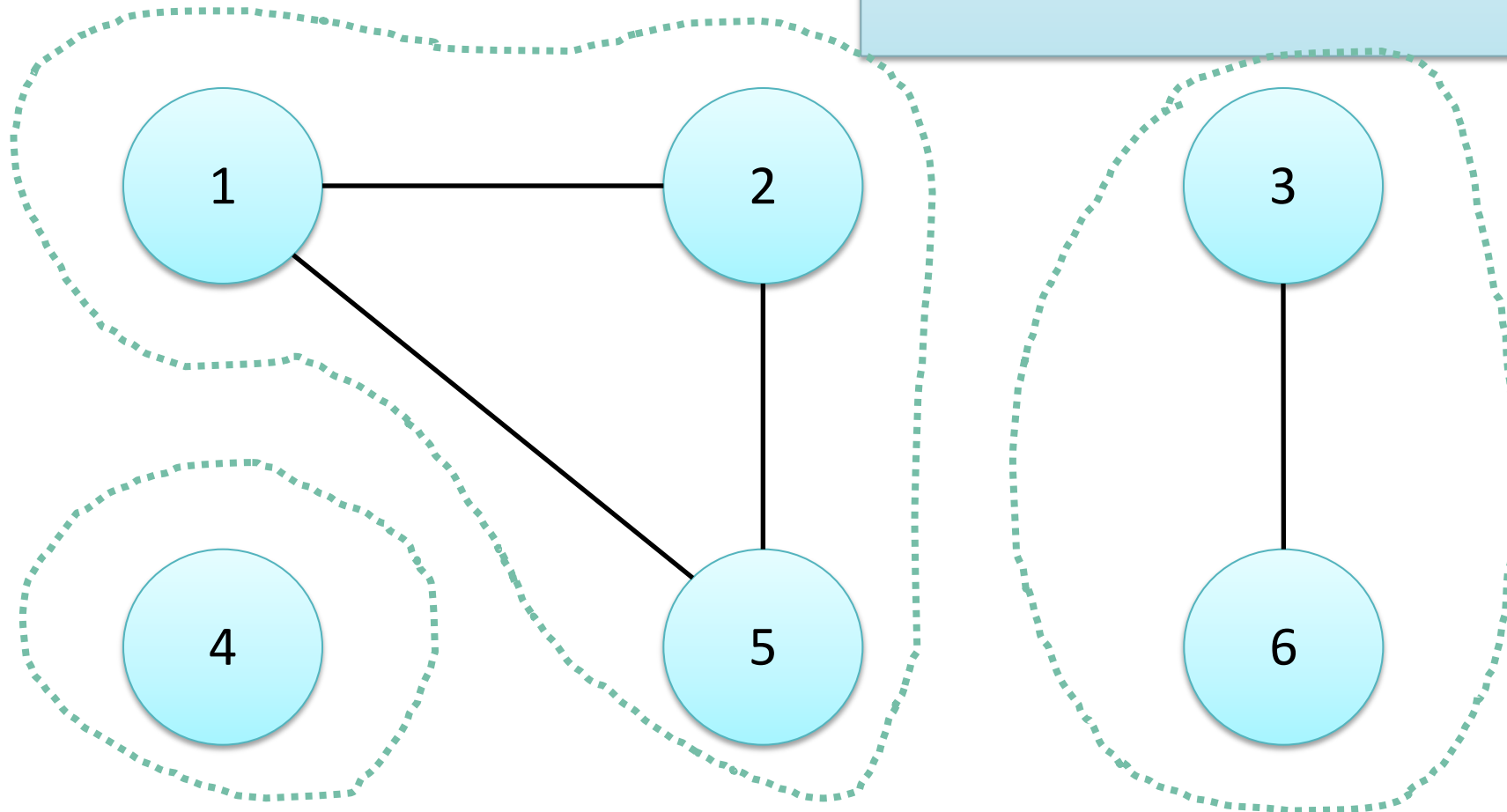


# Reachability (Undirected)

- An undirected graph is **connected** if, for every couple of vertices, there is a path connecting them
- The connected sub-graphs of maximum size are called **connected components**
- A connected graph has exactly one connected component

# Connected components

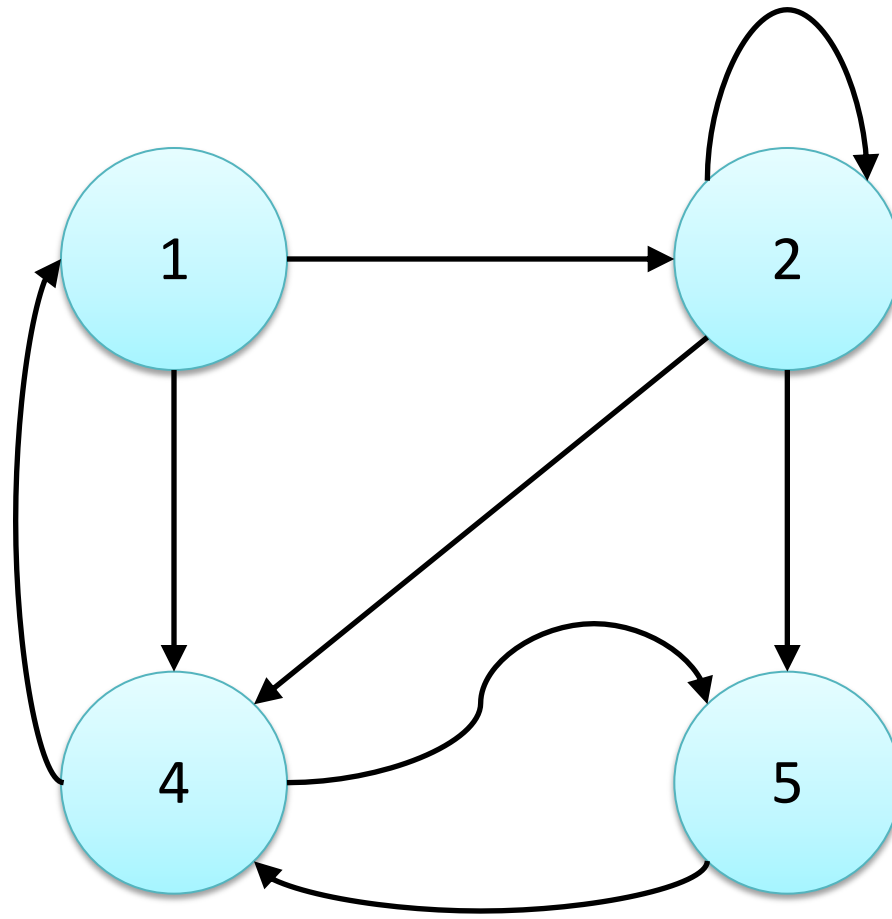
The graph is **not** connected.  
Connected components = 3  
 $\{4\}, \{1, 2, 5\}, \{3, 6\}$



# Reachability (Directed)

- A directed graph is **strongly connected** if, for every ordered pair of vertices  $(v, v')$ , there exists at least one path connecting  $v$  to  $v'$

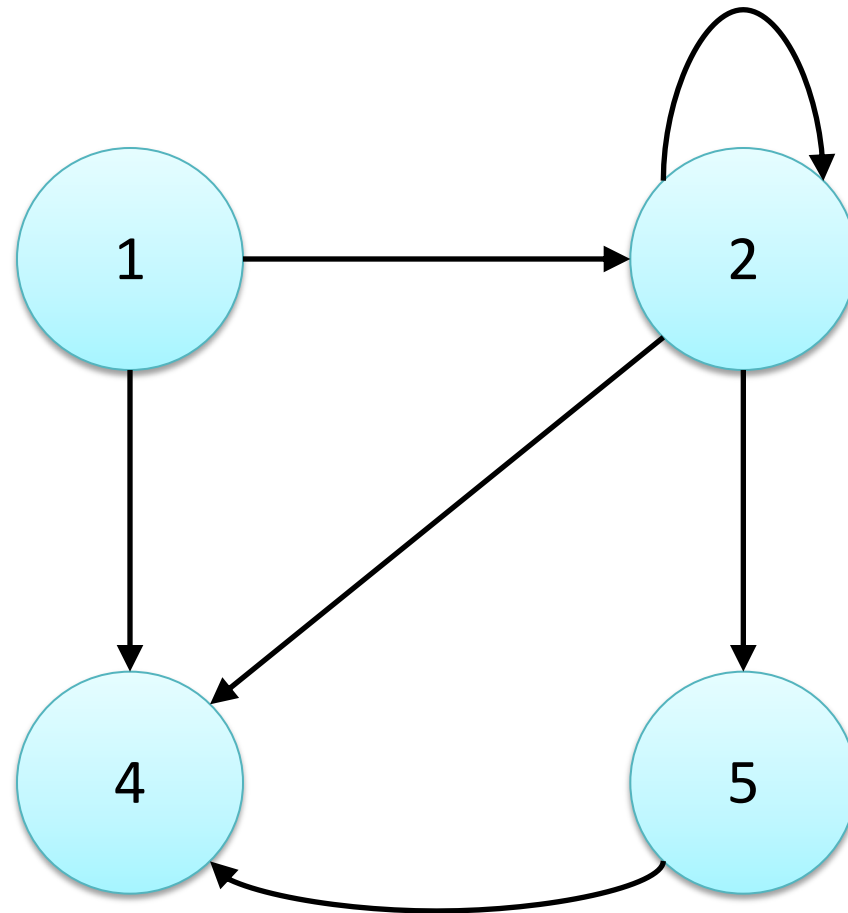
# Example



The graph is **strongly connected**



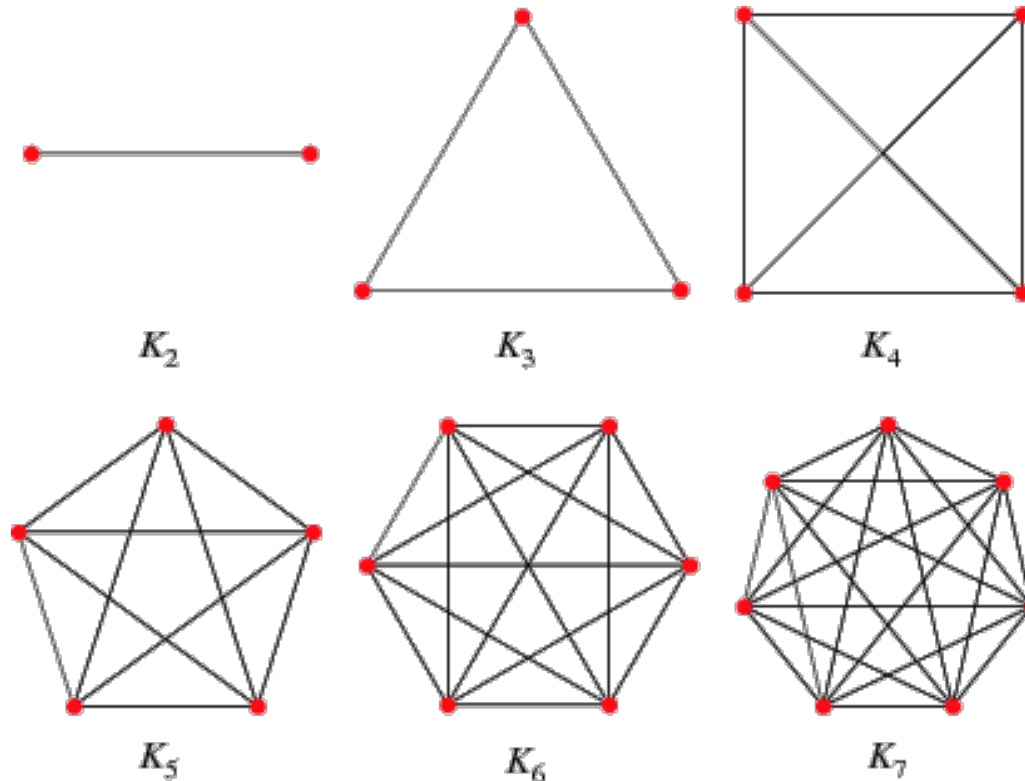
# Example



The graph is **not** strongly connected

# Complete graph

- A graph is complete if, for every pair of vertices, there is an edge connecting them (they are adjacent)
- Symbol:  $K_n$



# Complete graph: edges

- In a **complete** graph with  $n$  vertices, the number of **edges** is

	Directed	Undirected
No self loops	$n(n - 1)$	$\frac{n(n - 1)}{2}$
With self loops	$n^2$	$\frac{n(n + 1)}{2}$

# Density

- The density of a graph  $G=(V,E)$  is the ratio of the number of edges to the total number of possible edges

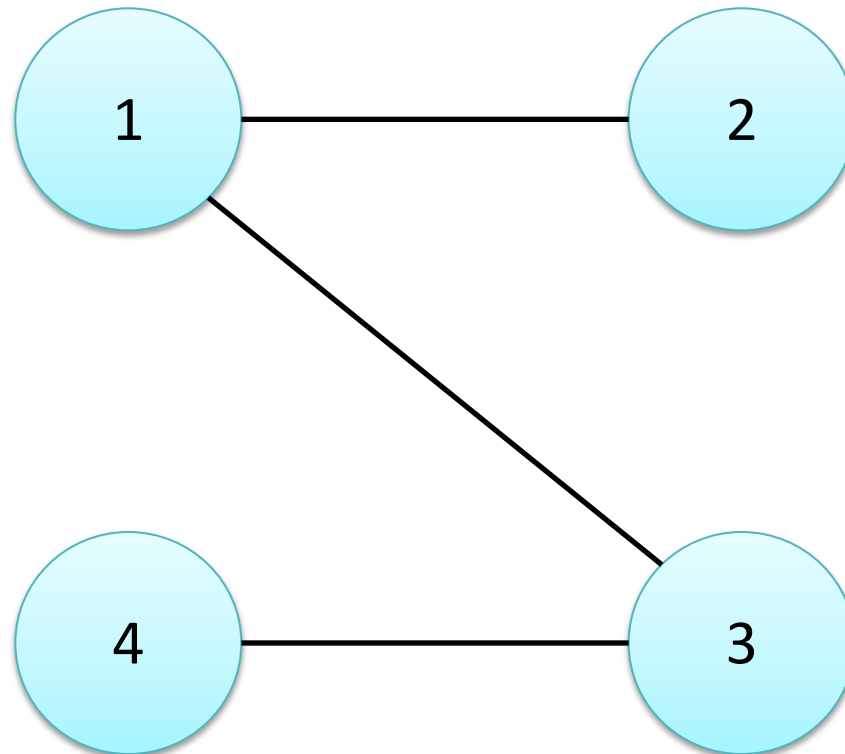
$$d = \frac{|E(G)|}{|E(K_{|V(G)|})|}$$

# Example

Density = 0.5

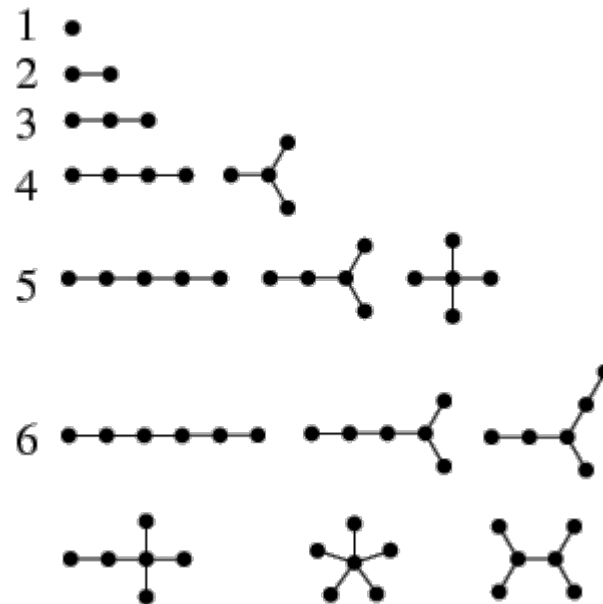
Existing: 3 edges

Total: 6 possible edges



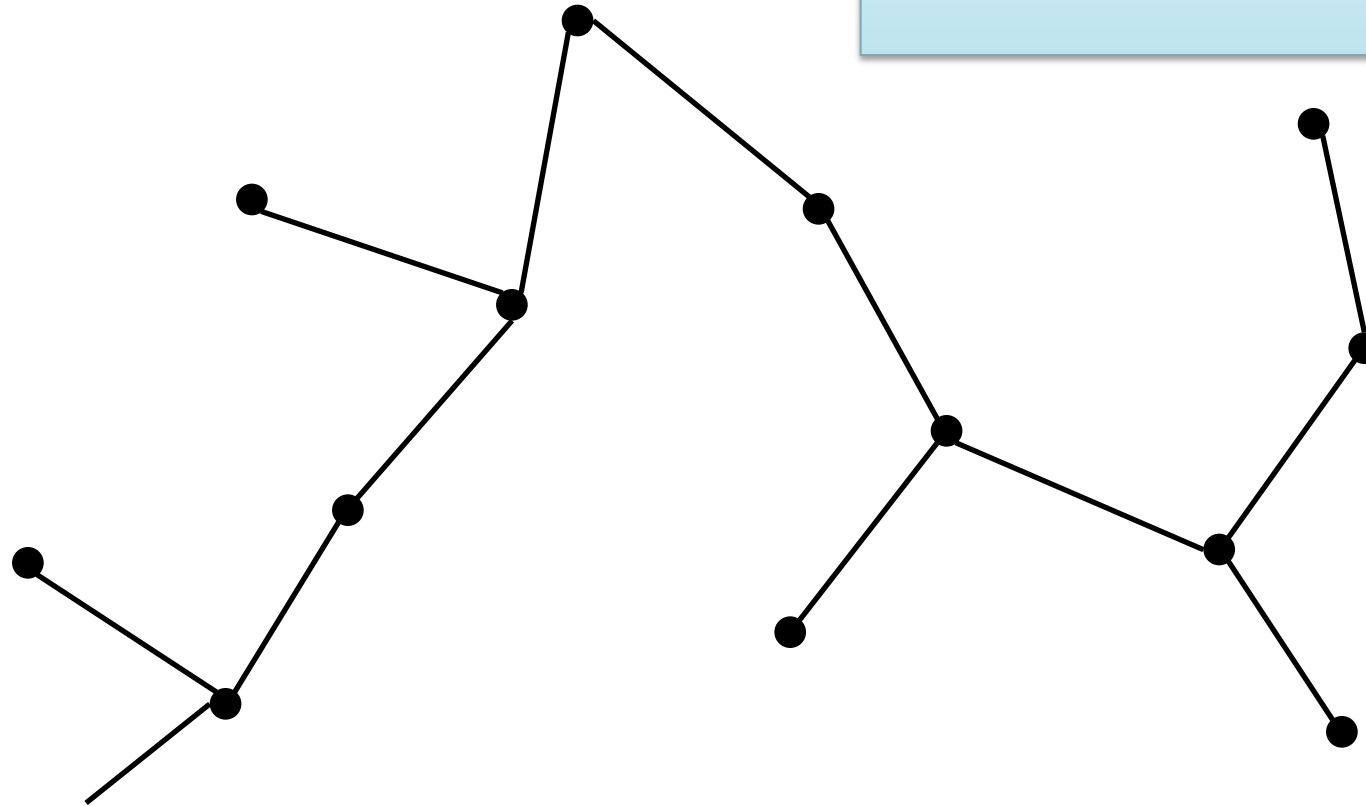
# Trees and Forests

- An undirected acyclic graph is called **forest**
- An undirected acyclic connected graph is called **tree**



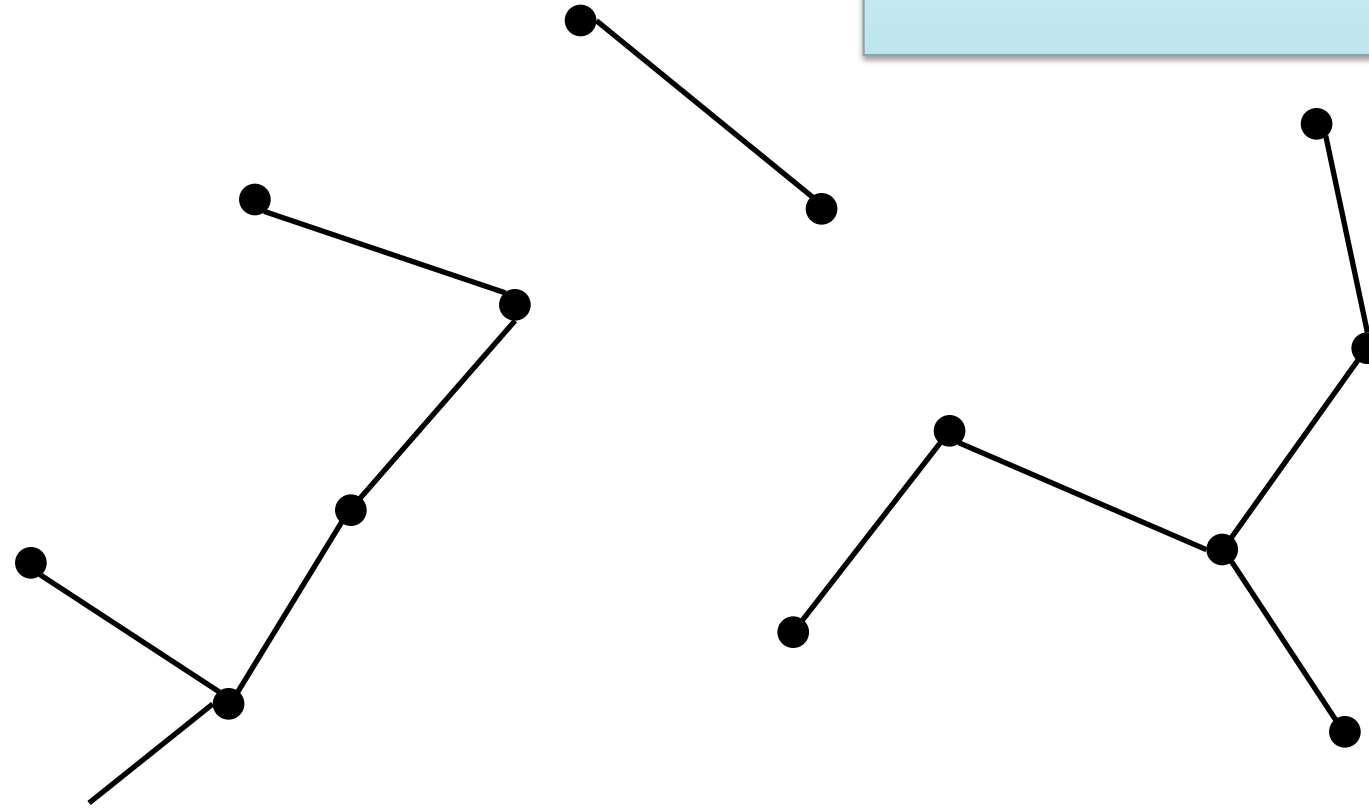
# Example

Tree



# Example

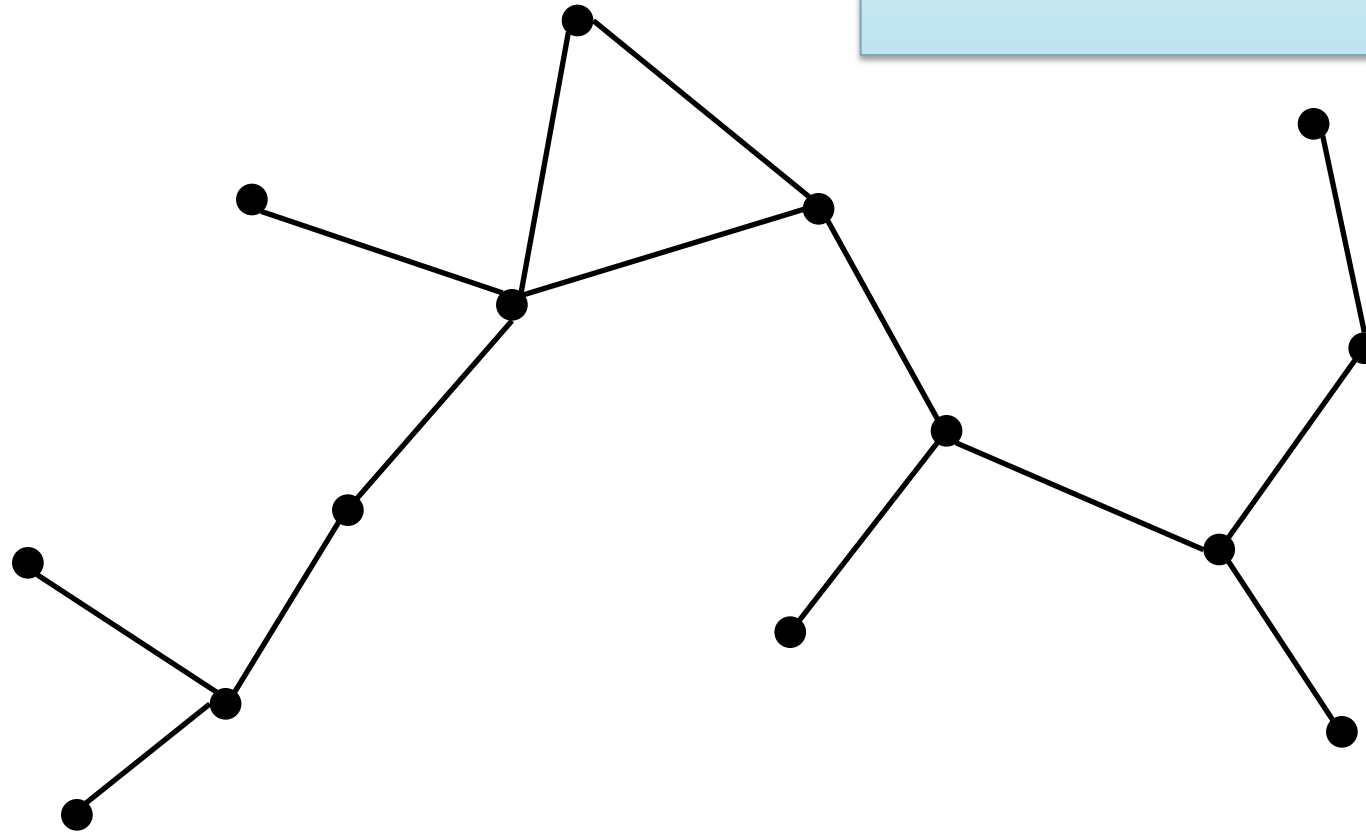
Forest





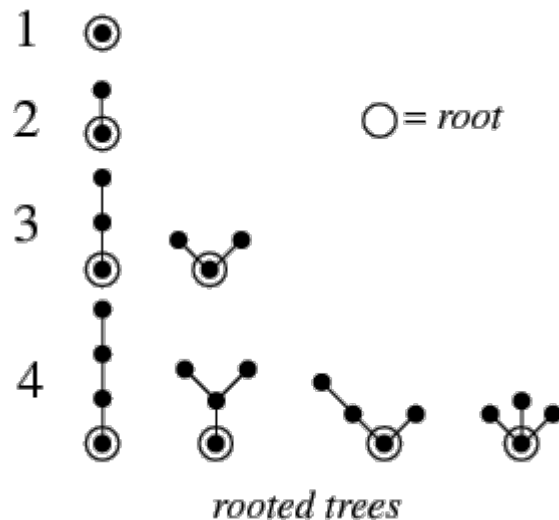
# Example

This is not a tree nor a forest  
(it contains a cycle)



# Rooted trees

- In a tree, a special node may be singled out
- This node is called the “**root**” of the tree
- Any node of a tree can be the root

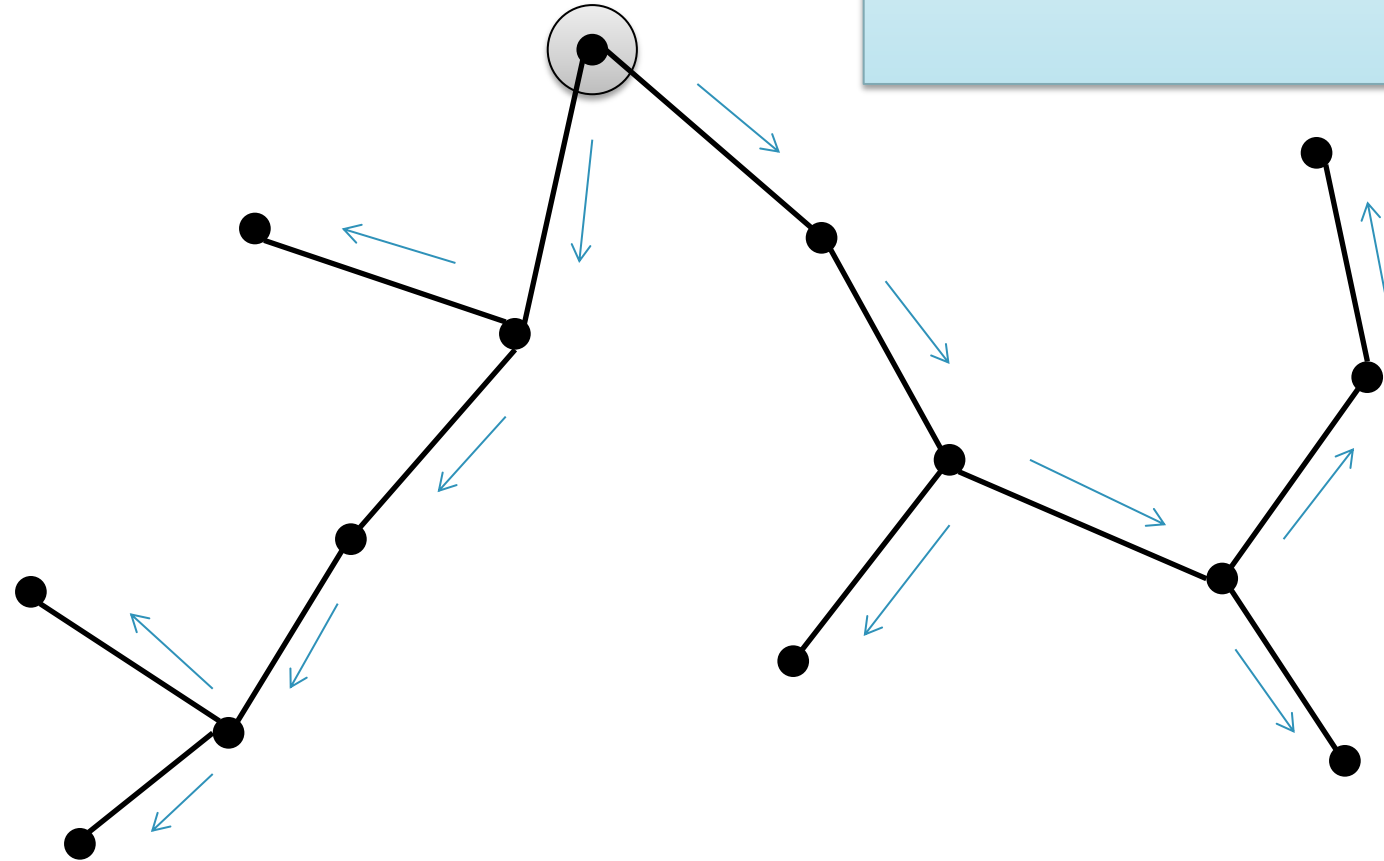


# Tree (implicit) ordering

- The root node of a tree **induces an ordering** of the nodes
- The root is the “ancestor” of all other nodes/vertices
  - “children” are “away from the root”
  - “parents” are “towards the root”
- The root is the only node without parents
- All other nodes have exactly one parent
- The furthestmost (children-of-children-of-children...) nodes are “leaves”

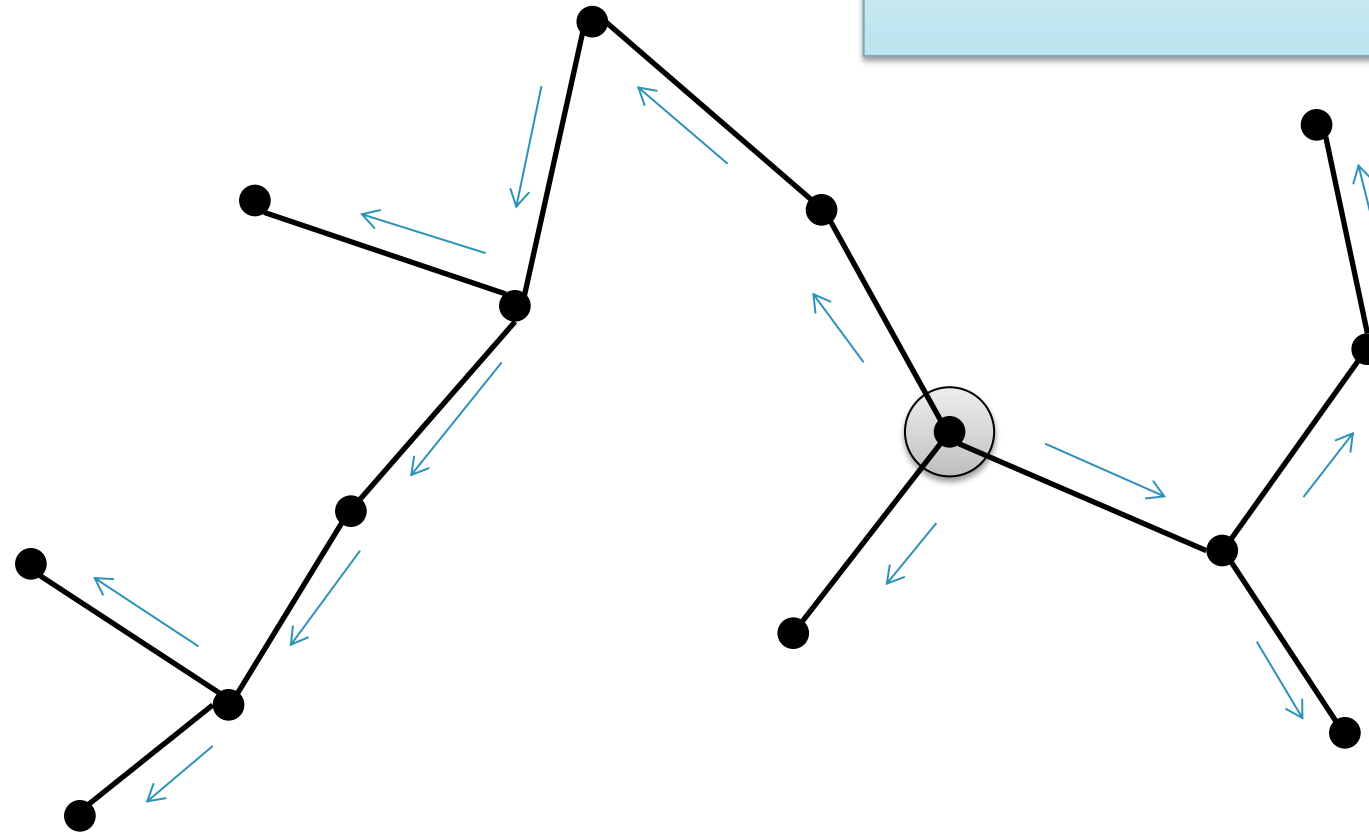
# Example

Rooted Tree



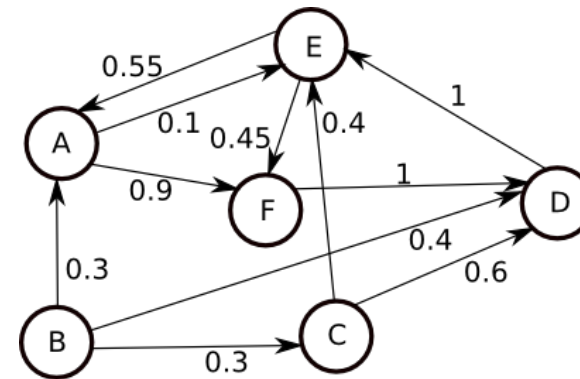
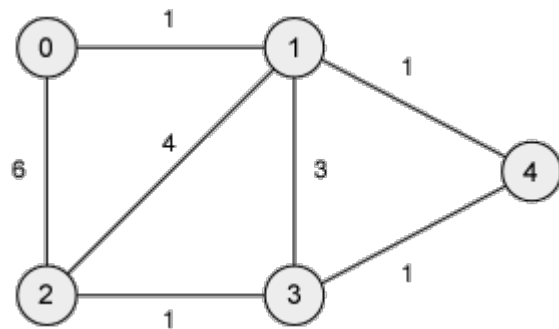
# Example

Rooted Tree

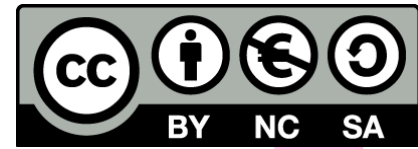


# Weighted graphs

- A weighted graph is a graph in which each branch (edge) is given a numerical weight.
- A weighted graph is therefore a special type of labeled graph in which the labels are numbers (which are usually taken to be positive).



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